

# 0x13 Ural Championship 2015

## Problem analysis

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# A. The first day of school

## Problem statement

- We have 3 days of school
- Each day has 4 time slots
- Each time slot can contain 1 lesson
- Each lesson has a name
- Name contains of up to 5 words, each up to 10 letters
- We need to print a timetable as an actual table

# A. The first day of school

## Problem solution

- Each cell has same width (10 characters)
- All cells in a row has equal height, which is equal to maximum height of cell in this row
- Height for single cell is calculated from the name of lesson in this cell
- The only thing we need is correct understanding of statement and thorough implementation

# B. Maths

## Problem statement

- We need to find a sequence  $a_i$  of length  $n$
- Constraint:  $\forall k \geq 2 : cnt(\sum_{i=1}^k a_i) = a_k$ , where  $cnt(x)$  is equal to amount of different positive divisors of  $x$

# B. Maths

## Problem solution

- Dynamic programming
- $dp_i$  is the maximum length of such sequence that can be finished with number  $i$
- $dp_i = dp_{i-cnt(i)} + 1$
- Answer is  $i$  for which  $dp_i = n$
- There is always an answer in  $[1; 2 \times 10^6]$  (proved by experiment)

# C. History

## Problem statement

- We have a time range  $[l; r]$ ,  $l$  and  $r$  are year numbers
- We need to calculate amount of years with all possible counts of Friday, 13th
- $l, r \leq 10^9$

# C. History

## Leap years rules

- Year is not leap, if it's number isn't divided by 4
- Year is not leap, if it's number is divided by 100, but isn't divided by 400
- All the other years are leap
- So the years  $x$  and  $x + 400$  are either both leap or both not leap

# C. History

## Finding a cycle

- Every 400 years contain the same amount of days in it (let's call it  $X$ )
- Every Every 2800 years contain the same amount of days  $7 \times X$ , and this amount is divided by 7
- So the years  $x$  and  $x + 2800$  always start from the same day of week



# C. History

## Solution

- We can calculate answer for the first 2800 years with day-by-day modeling
- Then we can multiply the answer on  $\frac{y-x+1}{2800}$
- And then we can model remaining years (there are less then 2800 of them)

# D. Chemistry

## Problem statement

- We have  $n$  tubes with 1 unit of liquid in each
- From two tubes with  $x$  and  $y$  ( $x \geq y$ ) units we can get two tubes with  $x - y$  and  $2 \times y$  units
- We need to get a tube with  $k$  units

# D. Chemistry

## Theorem

- For odd  $n$  we can get any  $x \in [0; n - 1]$
- For even  $n$  we can get any  $x \in [0; n - 2]$
- Algorithm will be derived from proof
- We'll use induction
- Base ( $n = 3, n = 4$ ) is obvious

# D. Chemistry

Induction step for even  $n$

- Consider even  $n$
- For  $n - 1$  we can get  $x \in [0; (n - 1) - 1 = n - 2]$

# D. Chemistry

## Observation

- Consider two tubes with  $a$  and  $b$ ,  $a > b$
- We can get two tubes with  $2 \times b = (2 \times b) \pmod n$  and  $a - b = a - n + a = 2 \times a - n = (2 \times a) \pmod n$
- So for one pour action we double amount in every tube by modulo 2
- So for  $t$  pour actions with  $a$  and  $b$  we get  $(2^t \times a) \pmod n$  and  $(2^t \times b) \pmod n$

## D. Chemistry

Induction step for odd  $n$

- Consider odd  $n$
- For  $n - 2$  we can get  $x \in [0; (n - 2) - 1 = n - 3]$
- Let's get two tubes with  $n - 4$  and 4 units
- Let's do  $\phi(n) - 2$  pour actions
- We'll get  $2^{\phi(n)} \bmod n = 1$  unit in second tube
- And we'll get  $n - 1$  unit in the first tube
- With one more pour action we'll get  $n - 2$  in the first tube

## D. Chemistry

### The rest of the proof (1)

- Theorem is almost proved
- Odd  $n$  can't be got (it should be the last pour action, but it's odd)
- For even  $n$   $n - 1$  can't be got (it should be got by pouring something from  $n$ , but in that case all other units are empty)
- For even  $n$   $n$  can't be got (proved on the next slide)

## D. Chemistry

### The rest of the proof (2)

- Consider going backwards
- We have one tube with  $n$  units and  $n - 1$  empty ones
- We can split any tube with even amount of units onto two parts
- If  $n$  is not some degree of 2, it will always be divided by some odd  $k \geq 3$
- Degree of 2 obviously can be got



# E. 3d-modeling

## Problem statement

- We have 2 straight lines
- We need to rotate one of them around some axis so that they became equal

# E. 3d-modeling

## Solution idea

- Let's set rotating angle to  $180^\circ$
- This way rotation is just reflecting from axis

# E. 3d-modeling

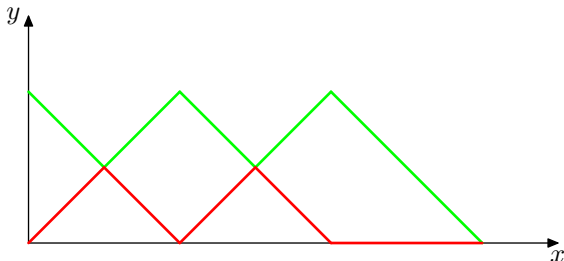
## Solution

- Let's say that  $\vec{a}$  and  $\vec{b}$  are direction vectors for our lines, and  $\vec{c}$  is direction vector for axis
- Let's say that  $X$  and  $Y$  are some pair of closest points on our lines
- If axis will go through  $\frac{X+Y}{2}$ , then  $Y$  will be reflected into  $X$
- If we set  $\vec{c} = \vec{a} + \vec{b}$ , then second line will be reflected into first line

# F. Physics

## Problem statement (1)

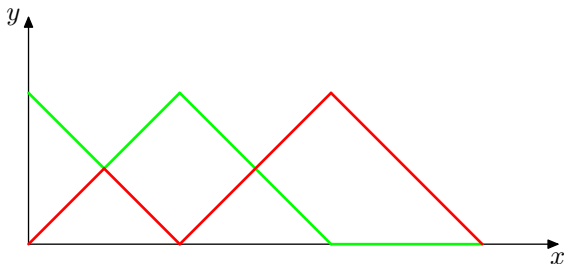
- We have two polygonal chains with some common points



# F. Physics

## Problem statement (2)

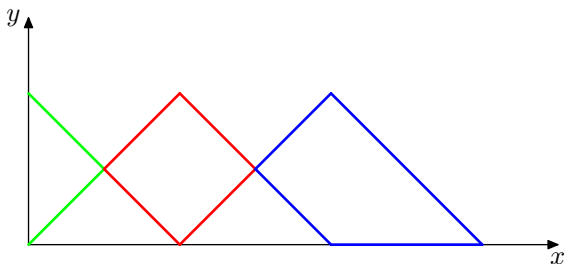
- We need to reassign some segments so that it would still be two polygonal chains
- Area under first chain should be equal to area under second



# F. Physics

## Solution idea (1)

- We have only 30 common points
- We can reassign only whole set of segments between two consequent common points
- Let's use common points to split our chains onto parts



# F. Physics

## Solution idea (2)

- For every part we'll calculate areas under it's top border and bottom border
- We have two sequences  $a_n$  and  $b_n$
- We need to find such set of indices  $S$  that
$$\sum_{i \in S} a_i + \sum_{i \notin S} b_i = \sum_{i \notin S} a_i + \sum_{i \in S} b_i$$

# F. Physics

## Solution idea (3)

- Let's split all indices onto two halves
- For every part we'll calculate all possible differences between  $\sum a$  and  $\sum b$
- For every possible difference in first half we'll try to find a matching difference in second half
- After that we'll only need to eliminate consequent triplets of points on same line



# G. Physical education

## Problem statement

- We sort all numbers from 1 to  $n$  ( $n \leq 10^9$ ) by sum of digits and then by value
- Need to find amount of numbers that didn't change their positions

# G. Physical education

## Problem solution

- Let's consider all numbers with equal digit sum
- They are sorted by their values
- We can precalculate set of positions they will take
- Difference between two consequent number is always greater, then 1
- So there is at most 1 number which hasn't change it's position
- It can be found by binary search
- Query *k-th number less then n with sum of digits equal to x* can be supported with dynamic programming

# H. Biology

## Problem statement

- Construct a planar graph with 16 vertices and  $\leq 300000$  simple cycles.

# H. Biology

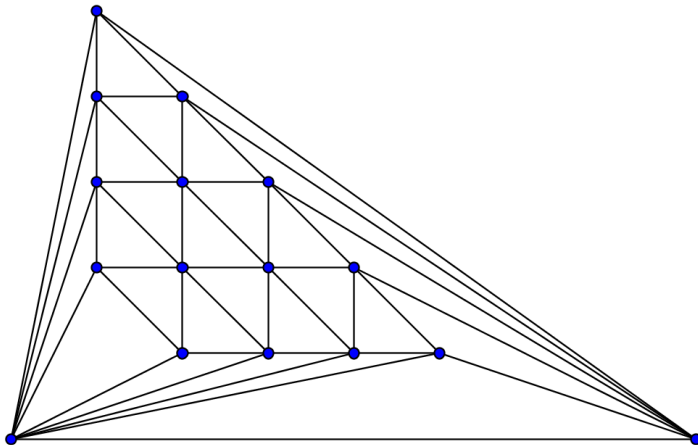
## Solution

- As for most constructive problems, write a local testing utility. Don't attempt to submit blindly! For this problem it's standard DP algorithm working in  $O(2^n n^2)$  time.
- After coding a local checker, try different approaches by hand, if they work, it's great!
- If that didn't help, try generating random tests. For this problem, generate random points and add random edges until you can't add any. It's also a good idea to randomly flip edges while the answer still increases.

# H. Biology

## Solution

- A picture of a solution, courtesy of nk.karpov @ Codeforces



# J. Urban geography

## Problem statement

- Given a weighted undirected graph
- Find the connected subgraph with minimum difference between maximum and minimum edge

# J. Urban geography

## Solution

- Sort all edges by weight, then run a binary search on problem's answer.
- How to check if the answer is  $\leq M$ ? Use two pointers technique to add and remove edges and check if graph becomes connected.
- It's hard to support all the above operations online, let's do this offline! Dynamic connectivity offline algorithm can be implemented to work in  $O(M \log^2 N)$ .
- Summary: binary search + offline dynamic connectivity gives  $O(M \log^2 N \log C)$  algorithm.